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DB=USPT; PLUR=YES; OP=ADJ

<u>L15</u>	(predecessor or forefather or ancestor) same access\$3 near5 right	3	<u>L15</u>
<u>L14</u>	database same (predecessor or forefather or ancestor) same access\$3 near5 right	0	<u>L14</u>
<u>L13</u>	L12 and (parent or child or ancestor or predecessor)	12	<u>L13</u>
<u>L12</u>	database near5 access\$3 right	51	<u>L12</u>
<u>L11</u>	l10 and ancestor	3	<u>L11</u>
<u>L10</u>	5941947[uref]	53	<u>L10</u>
<u>L9</u>	5941947.pn.	1	<u>L9</u>
<u>L8</u>	ancestor group\$3	4	<u>L8</u>
<u>L7</u>	(user or client) same ancestor group	0	<u>L7</u>
<u>L6</u>	(user or client)near5 group\$3 same ancestor group	0	<u>L6</u>
<u>L5</u>	(assign\$3 or classif\$6) same (user or client) same group\$3 same ancestor group	0	<u>L5</u>
<u>L4</u>	(assign\$3 or classif\$6) same (user or client) near5 group\$3 same ancestor group same access\$3	0	<u>L4</u>
<u>L3</u>	l1 same database	14	<u>L3</u>
<u>L2</u>	l1 same hierarchical database	1	<u>L2</u>
<u>L1</u>	(node or group or client or user) same ancestor (node or group or client or user)	204	<u>L1</u>

END OF SEARCH HISTORY

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L8: Entry 2 of 4

File: USPT

Sep 10, 1996

US-PAT-NO: 5555375

DOCUMENT-IDENTIFIER: US 5555375 A

TITLE: Method and apparatus for network computer systems management group administration

DATE-ISSUED: September 10, 1996

INVENTOR-INFORMATION:

NAME	CITY	STATE	ZIP CODE	COUNTRY
Sudama; Ram	Hudson	MA		
Griffin; David M.	Maynard	MA		
Johnson; Brad	Westerly	RI		
Sealy; Dexter	Boston	MA		
Shelhamer; James	Maynard	MA		
Tallman; Owen H.	Lunenburg	MA		

US-CL-CURRENT: 709/226; 707/10

ABSTRACT:

The following is a method and apparatus for administering an operation specified for performance on a set of independently managed hosts. The operation is received initially by one of a plurality of management servers in a managed host system. The operation, specifying a group object, is transferred to the management server designated by the system to administer the operation specified on the group object. The designated management server thereafter decomposes the group object into constituent objects which may be host objects or additional group objects. Next, the locally administered objects are scheduled for execution on the hosts administered by the designated management server. The non-locally administered objects are transferred to the management servers identified in a database for administering the objects. After executing the operation on the host objects, the host objects and group objects return status information back to the designated management server. The status information is then transmitted to the management server that initially received the operation.

8 Claims, 4 Drawing figures

Exemplary Claim Number: 1

Number of Drawing Sheets: 4

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L13: Entry 1 of 12

File: USPT

Feb 4, 2003

US-PAT-NO: 6516317

DOCUMENT-IDENTIFIER: US 6516317 B1

TITLE: Method and apparatus for facilitating compartmentalized database user management

DATE-ISSUED: February 4, 2003

INVENTOR-INFORMATION:

NAME	CITY	STATE	ZIP CODE	COUNTRY
Samar; Vipin	Cupertino	CA		
Lei; Chon Hei	Alameda	CA		
Keefe; Thomas F.	San Mateo	CA		

US-CL-CURRENT: 707/9

ABSTRACT:

One embodiment of the present invention provides a system that facilitates compartmentalized user management in a database system. This database system is compartmentalized into a plurality of domains that are insulated from each other, so that a given user who has access to data within an associated domain does not have access to data in other domains. Upon receiving a request from a database administrator to perform an operation on a user within the database system, the system identifies a domain that the user is associated with in the database system. Next, the system determines whether the database administrator is authorized to modify users associated with the domain. If so, the system performs the operation by modifying an entry for the user within a user table in the database system that contains an entry for each user of the database system. In one embodiment of the present invention, the operation on the user can include: creating the user within the database system; altering the domain that the user is associated with in the database system; and dropping the user from the database system.

27 Claims, 12 Drawing figures

Exemplary Claim Number: 1

Number of Drawing Sheets: 8

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L4: Entry 1 of 4

File: USPT

Dec 2, 2003

US-PAT-NO: 6658413

DOCUMENT-IDENTIFIER: US 6658413 B1

TITLE: Multidimensional database system with intermediate lockable intersections

DATE-ISSUED: December 2, 2003

INVENTOR-INFORMATION:

NAME	CITY	STATE	ZIP CODE	COUNTRY
Reddy; Venugopal P.	Madison	NJ		
Matharu; Harmindar S.	Lake Hiawatha	NJ		

US-CL-CURRENT: 707/8; 707/101, 707/2, 707/9

ABSTRACT:

A multi-dimensional database includes hierarchical levels. The data is hierarchical, with aggregatable data measures, so that changes in a first intersection will cause changes in other intersections for the aggregated measure. Selected intersections can be locked, so that changes in aggregatable data are restricted. Constrained sets corresponding to the locked intersections limit the scope of changes which may be made, and ensure a high level of efficiency.

3 Claims, 11 Drawing figures

Exemplary Claim Number: 1

Number of Drawing Sheets: 5

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L4: Entry 1 of 4

File: USPT

Dec 2, 2003

DOCUMENT-IDENTIFIER: US 6658413 B1

TITLE: Multidimensional database system with intermediate lockable intersections

Abstract Text (1):

A multi-dimensional database includes hierarchical levels. The data is hierarchical, with aggregatable data measures, so that changes in a first intersection will cause changes in other intersections for the aggregated measure. Selected intersections can be locked, so that changes in aggregatable data are restricted. Constrained sets corresponding to the locked intersections limit the scope of changes which may be made, and ensure a high level of efficiency.

Brief Summary Text (11):

In accordance with the present invention, a multi-dimensional database includes hierarchical levels. The data is hierarchical, with aggregatable data measures, so that changes in a first intersection will cause changes in other intersections for the aggregated measure. Selected intersections can be locked, so that changes in aggregatable data are restricted. Constrained sets corresponding to the locked intersections limit the scope of changes which may be made, and ensure a high level of efficiency.

Detailed Description Text (122):

A custom level of a level X within an ancestor Y is defined by merging sub-groups of members of X that are related to a single member of Y. This could be performed for multiple members of Y. The members of new level Z thus formed are assigned new codes and descriptions and the level has the following relations:

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L4: Entry 2 of 4

File: USPT

Feb 5, 2002

DOCUMENT-IDENTIFIER: US 6345277 B1

TITLE: Method and apparatus for using an information model to organize an information repository into an extensible hierarchy of information

Detailed Description Text (12):

In presenting a hierarchical and logical view of information in a database, an information model allows a user to see two things. First, a hierarchy gives a user the ability to see how data in a database is organized and the relationship between data, i.e. fields, in the database. When a derived hierarchy is created in a preferred embodiment, derived containers are displayed such that hierarchical organization is conveyed to the user. In other words, a derived container can convey a category of information in the database as defined by its corresponding value-defined container definition node. A category of information corresponds to a field from the database from which its selection criteria is based, and the field can be used to create a label of the derived container. In a first embodiment, then, each derived container represents a category of information. Furthermore, if a derived container corresponds to a value-defined container definition node in the second level of information in the information model, a user interface can convey this sense of hierarchy by indenting the label of the derived container in a list user interface, or presenting a secondary window in a graphical window interface, for example.

Detailed Description Text (56):

Can all values of the field attribute be represented so that top level hierarchical relationships are maintained? In other words, every derived container in the same container group should be logical children/descendants of their parents/ancestors.

Detailed Description Text (66):

Third-level derived containers Development 2306 and Marketing 2320 are based on the same field (their selection criteria is field:department as will be illustrated). Furthermore, if all values of the field department in the database of FIG. 22 were represented, top level hierarchical relationships would be maintained since other values, SHP/REC and ADMIN are logical descendants of derived container Dept. List 2304. However, since not all derived containers can logically comprise the same children, this container group is not appropriately extensible. For instance, derived containers Office 2308 and Parts 2310, which are child derived containers of Development 2306, would not be logical children of SHP/REC because a shipping and receiving department would probably not comprise office type assets.